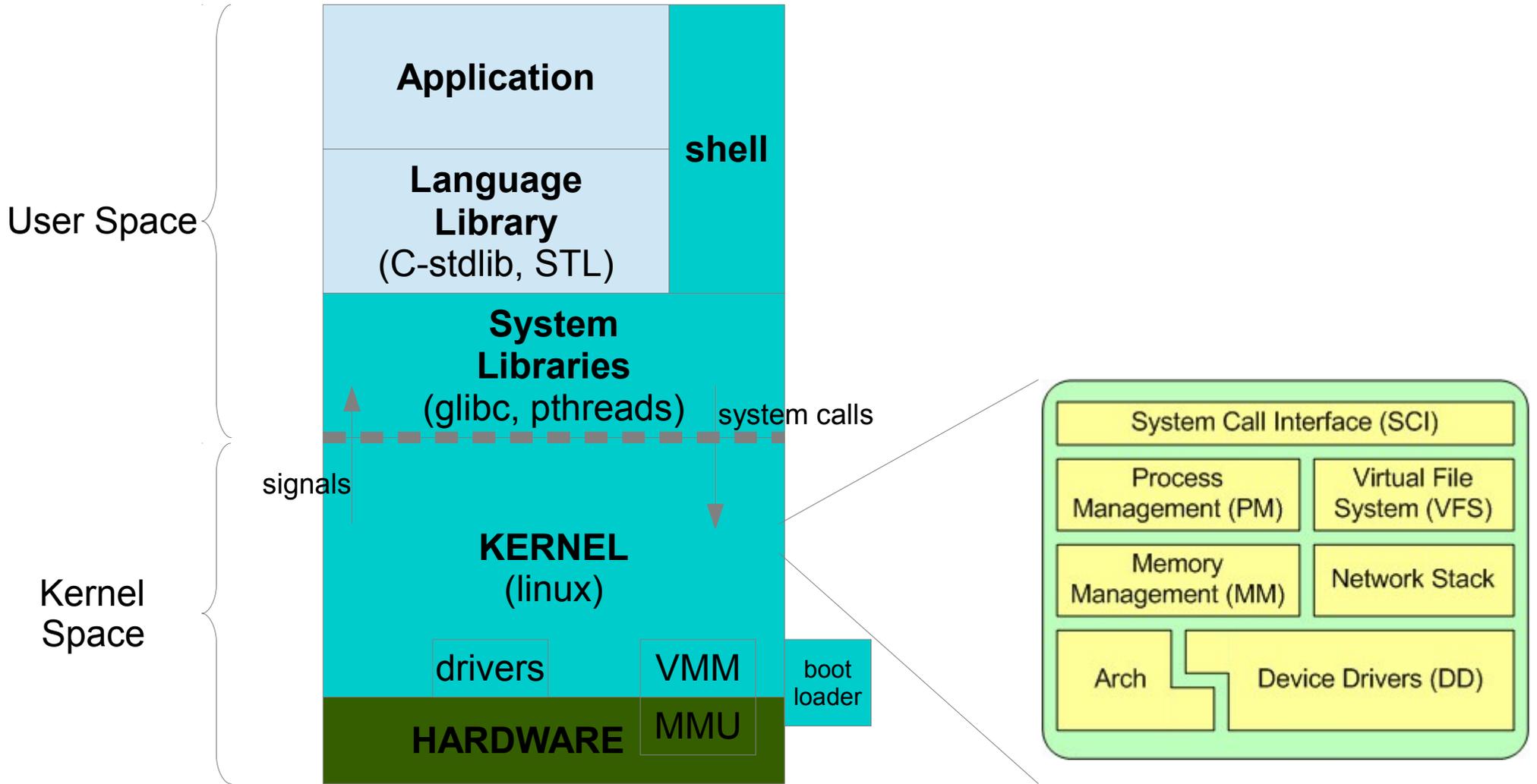
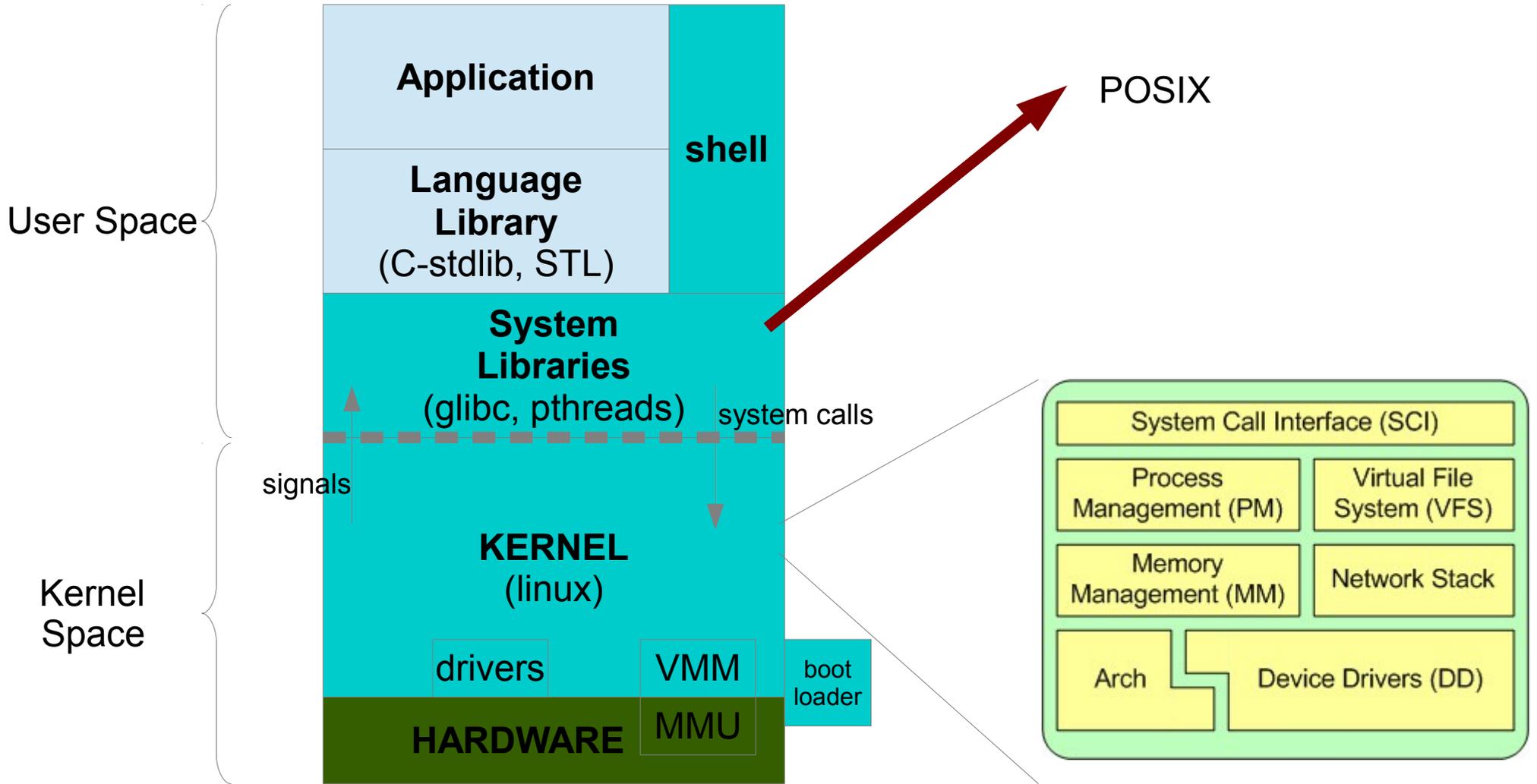


GCC internals intro y optimizationes

Daniel Gutson

Contexto: stack general





Toolchain

- Conjunto de herramientas

Toolchain

- Conjunto de herramientas
 - No todas del mismo fabricante

Toolchain

- Conjunto de herramientas
 - No todas del mismo fabricante
 - Algunas lenguaje-específicas

Toolchain

- Conjunto de herramientas
 - No todas del mismo fabricante
 - Algunas lenguaje-específicas
- Construye binarios

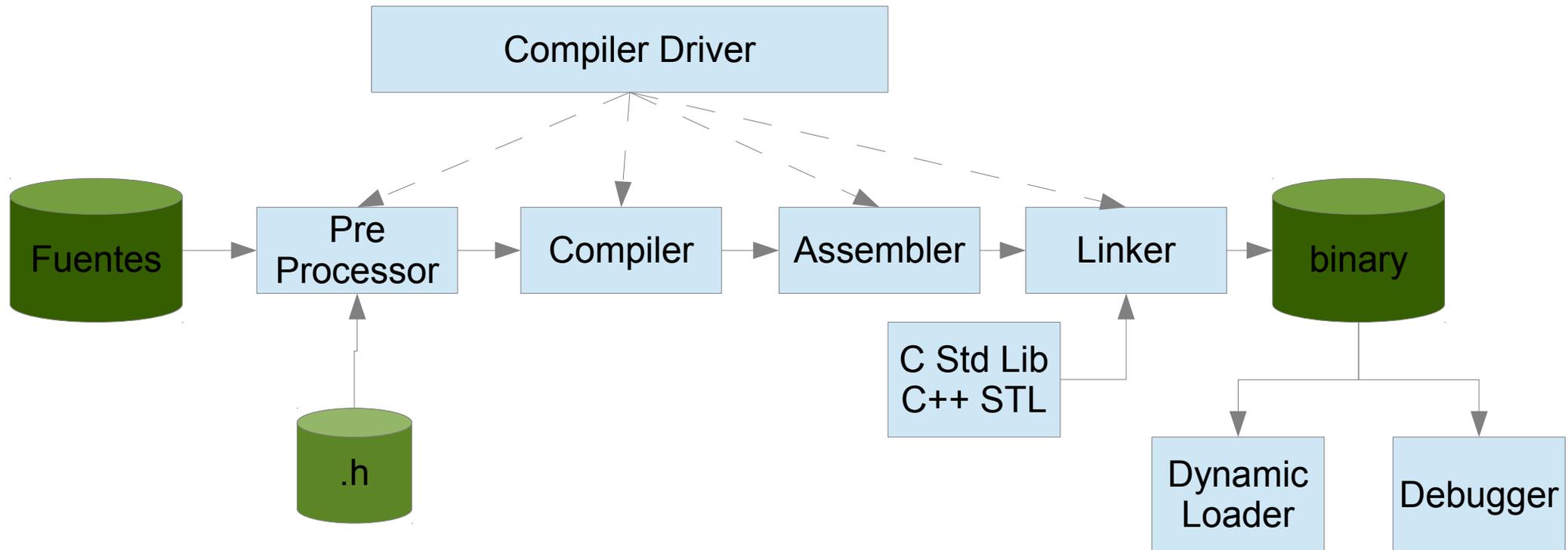
Toolchain

- Conjunto de herramientas
 - No todas del mismo fabricante
 - Algunas lenguaje-específicas
- Construye binarios
 - Responden a una ABI

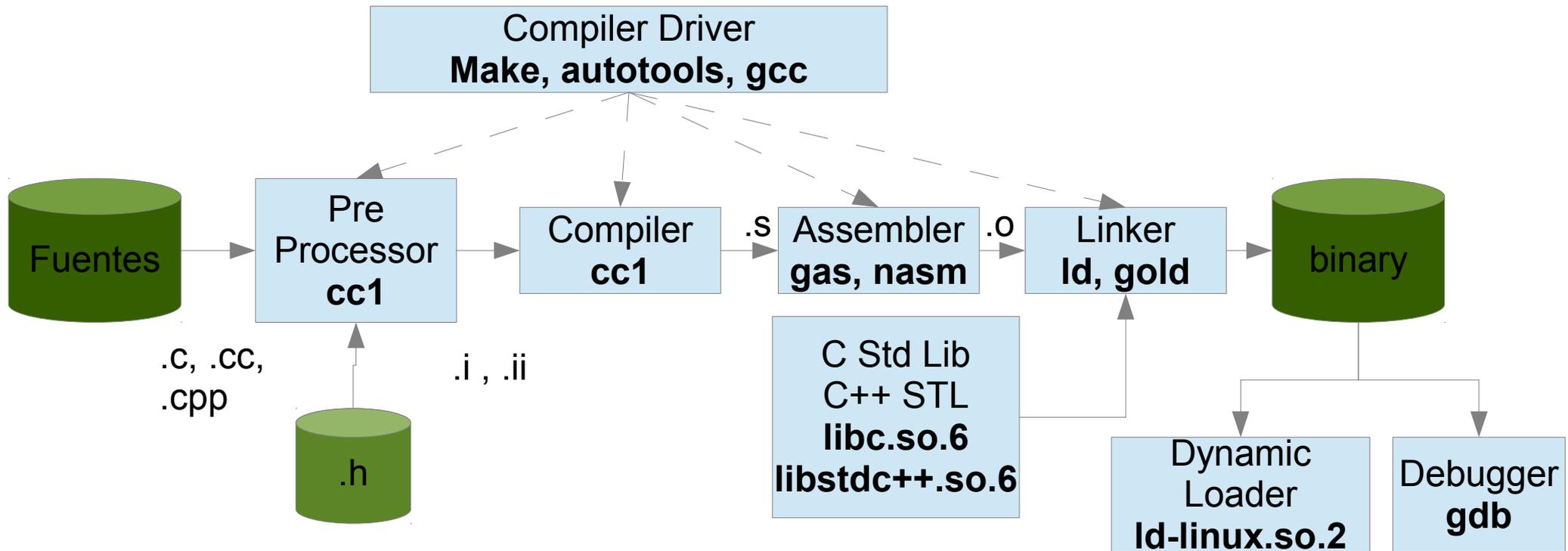
Toolchain

- Conjunto de herramientas
 - No todas del mismo fabricante
 - Algunas lenguaje-específicas
- Construye binarios
 - Responden a una ABI
- Características
 - Host arch.
 - Target arch.

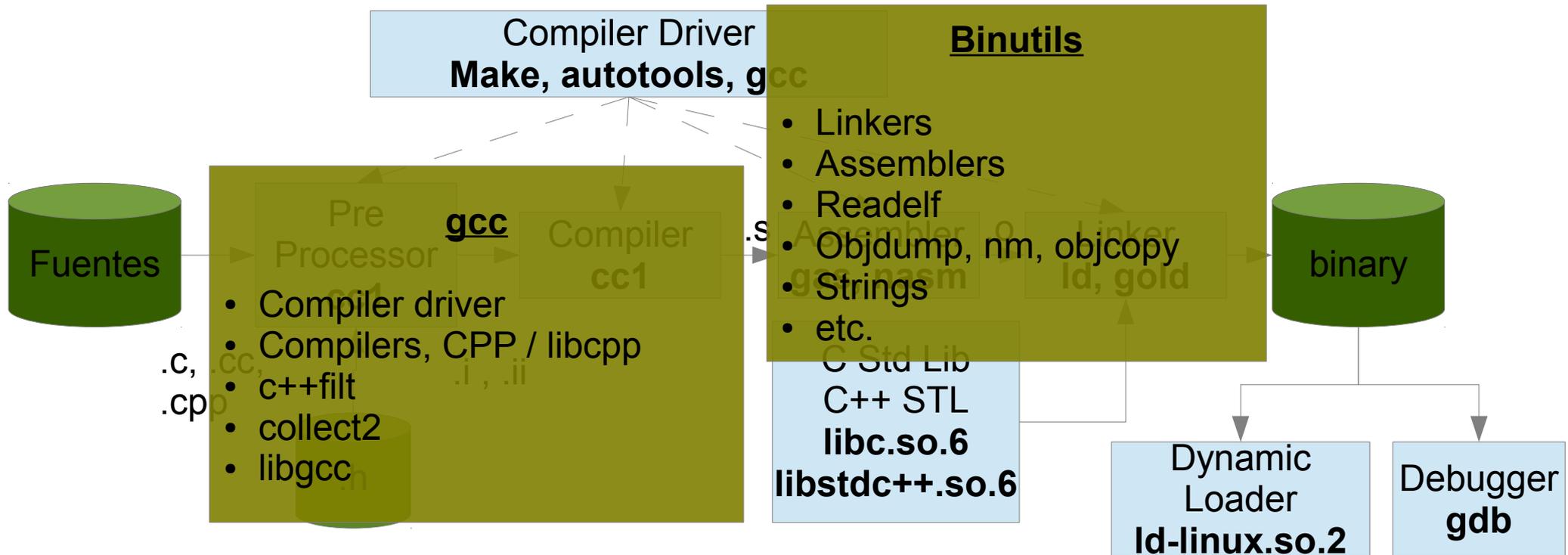
C & C++ Toolchain General



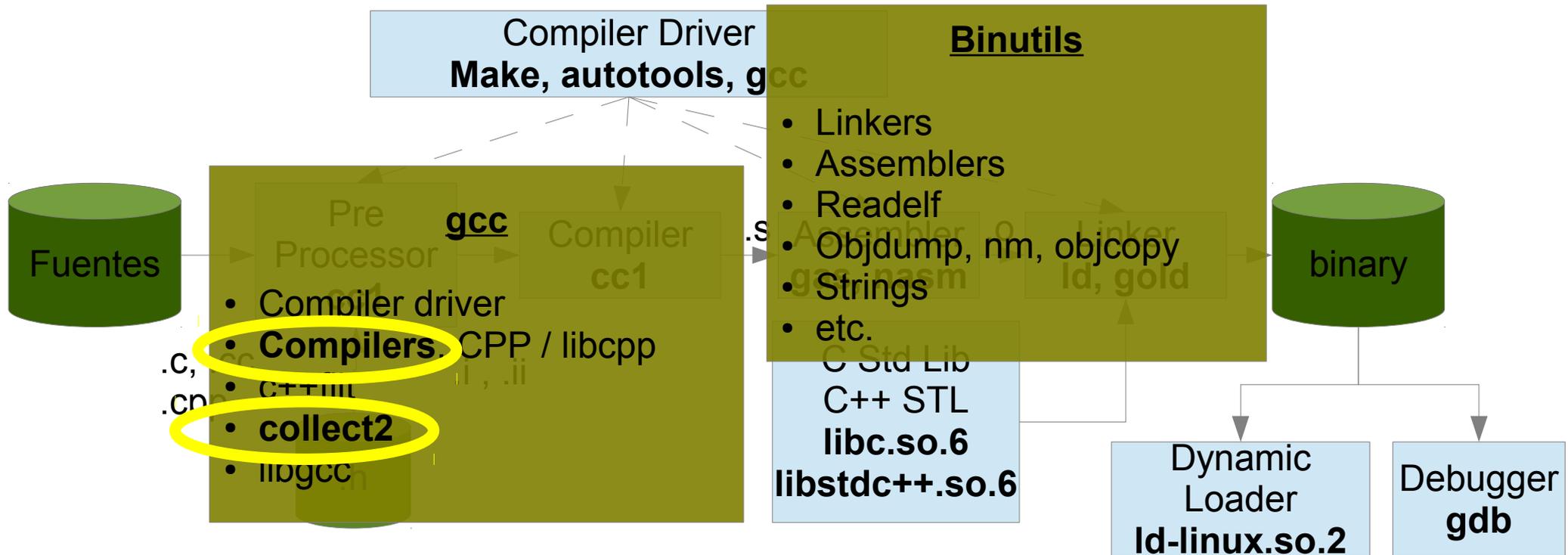
C & C++ Toolchain GNU



C & C++ Toolchain GNU

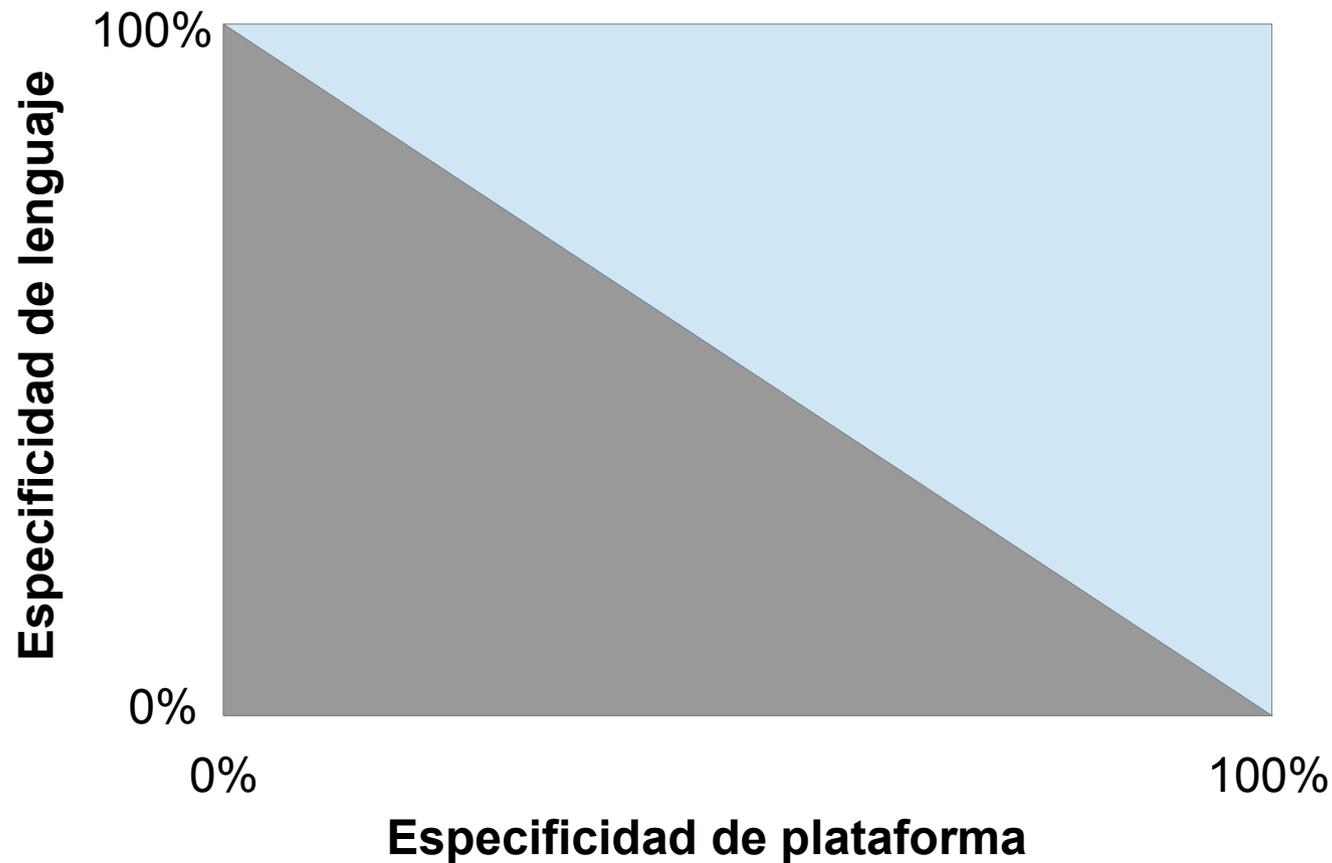


C & C++ Toolchain GNU



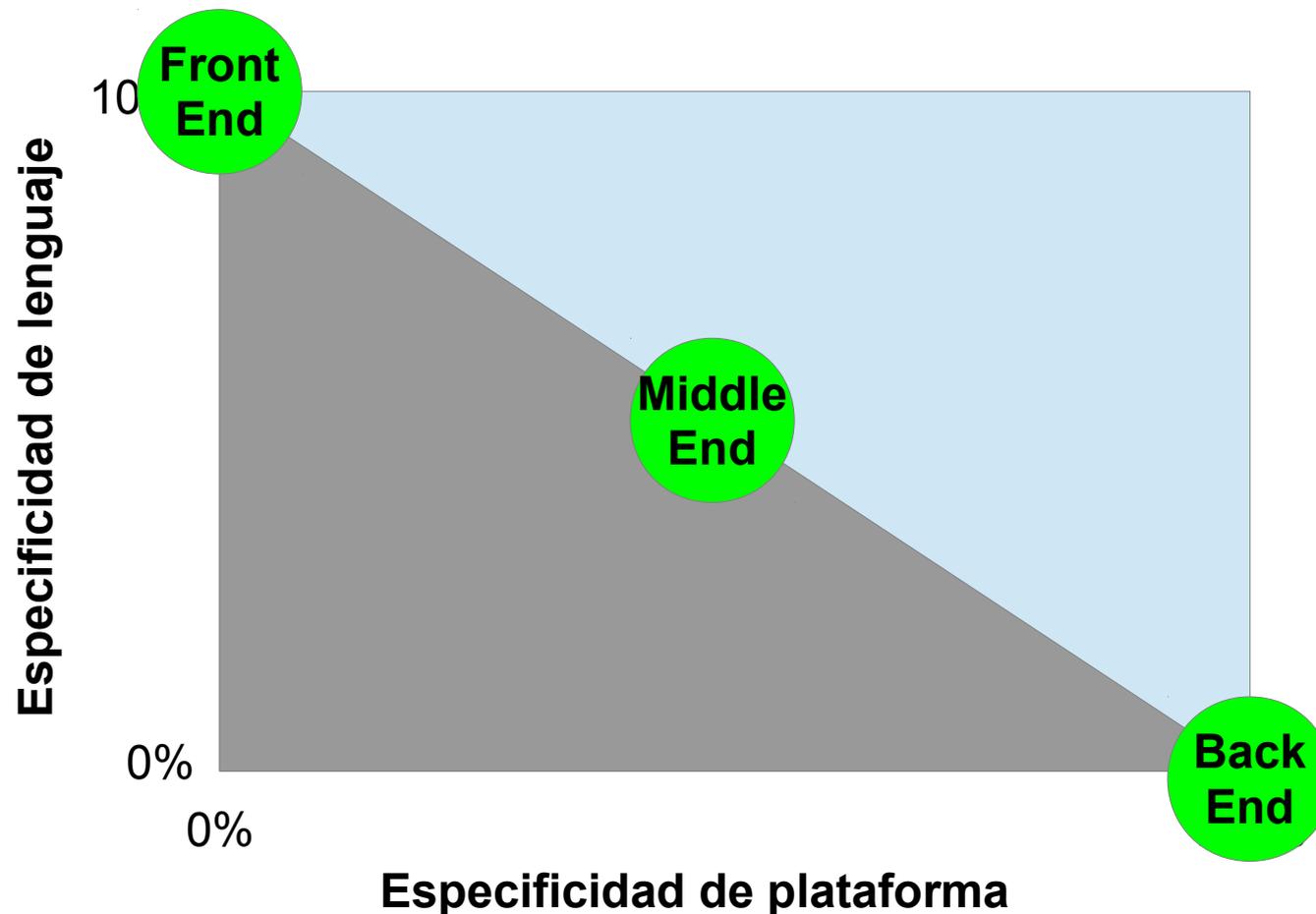
La colección de compiladores: GCC

- Comparten arquitectura común



La colección de compiladores: GCC

- Comparten arquitectura común



La colección de compiladores: GCC

- Comparten arquitectura común
- Comparten karma común
 - La antipatía de los que hacen el linker

Opciones de GCC

- **-fxxxxxxx** **F**ront-end específicas
- **-mxxxxxxx** **M**achine (back-end específicas)
- **-Wxxxxxxx** **W**arnings (todos-los-end)
- **-Oxxxxxxx** **O**ptimization (todos-los-end)

Opciones de GCC

- **-fxxxxxxx** **F**ront-end específicas
- **-mxxxxxxx** **M**achine (back-end específicas)
 - **-marchxx** **A**rchitecture
 - **-mcpuxxx** **C**PU-específicas
- **-Wxxxxxxx** **W**arnings (todos-los-end)
- **-Oxxxxxxx** **O**ptimization (todos-los-end)

Optimizaciones en GCC

- Optimizaciones en Front-End
 - Reconocen language-idioms
- Optimizaciones en Middle-end
 - Transforman loops (ej Graphite)
- Optimizaciones en Back-end
 - Selección de instrucciones apropiadas
 - Instruction scheduling
- Optimización por vectorización
- WPO (optimizaciones que deberían ser LTO)

Optimizaciones del Front-end C++

- Inlining C++-específico
- TMP
- RVO
- Copy Elision
- Eliminación de código muerto

Don't reach backend: TMP

- **Declarative**

- Integer Arithmetic Calculi
 - Uses enums
- Type calculi

- **Imperative**

- Float Arithmetics
- Expression Templates
- Loop Transformations
- AOP

=> All use & depend on inlining

-O3, -ftemplate-depth, cache, etc.

TMP Declarative – Example: Factorial

```
template <unsigned int N>
struct Facto
{
    enum { result = N * Facto<N-1>::result };
};

template <>
struct Facto<0>
{
    enum { result = 1 };
};

int main()
{
    return Facto<4>::result;
}
```

```
08048494 <main>:
8048494: 55          push  %ebp
8048495: 89 e5      mov   %esp,%ebp
8048497: b8 18 00 00 00 mov   $0x18,%eax
804849c: 5d        pop   %ebp
804849d: c3        ret
```

TMP Imperative – Example: PI

```
template <unsigned int N> struct PI {
    static float sign()    {
        return 1 - (2*(signed(N) % 2));
    }

    static float result() {
        return PI<N-1>::result() + sign() * 4.0f / ((2*N)+1);
    }
};

template <> struct PI<0>
{
    static float result() { return 4.0f; }
};

int main() {
    cout << PI<50>::result() << endl;
    return 0;
}
```

TMP Imperative – Example: PI

```
template <unsigned int N> struct PI {
    static float sign()    {
        return 1 - (2*(signed(N) % 2));
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template <> struct PI<0>
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};

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    cout << PI<50>::result() << endl;
    return 0;
}
```



TMP Imperative – Example: PI

```
template <unsigned int N> struct PI {  
    static float sign()    {  
        return 1 - (2*(signed(N) % 2));  
    }  
  
    static float result() {  
        return PI<N-1>::result() + sign() * 4.0f / ((2*N)+1);  
    }  
};
```

```
te  
{  
};  
in  
}  
08048790 <main>:  
8048790: 55          push  %ebp  
8048791: 89 e5      mov   %esp,%ebp  
8048793: 83 e4 f0   and  $0xffffffff0,%esp  
8048796: 56          push  %esi  
8048797: 53          push  %ebx  
8048798: 83 ec 18   sub  $0x18,%esp  
804879b: d9 05 d0 88 04 08 flds 0x80488d0  
80487a1: dd 5c 24 04 fstpl 0x4(%esp)  
80487a5: c7 04 24 40 a0 04 08 movl $0x804a040,(%esp)  
80487ac: e8 bb fe ff ff      call 804866c <std::basic_ostream<char,  
std::char_traits<char> >& std::basic_ostream<char, std::char_traits<char>  
>::_M_insert<double>(double)@plt>
```

Special Methods – RVO

```
struct S
{
    ...
};

S f()
{
    S ret;
    // manipulate ret
    return ret;
}

void g()
{
    S s = f();
}
```

Special Methods – RVO

```
struct S
{
    ...
};

S f()
{
    S ret;
    // manipulate ret
    return ret;
}

void g()
{
    S s = f();
}
```

The diagram illustrates Return Value Optimization (RVO) in C++. It shows a function `f()` that returns a struct `S`. The return value is represented by a yellow circle around `S ret;`. A green arrow points from this circle to another yellow circle around `S s = f();` in the function `g()`, indicating that the compiler can optimize the return value by constructing it directly in the caller's scope, avoiding a copy.

Special Methods – RVO

```
struct S
{
    S() { cout << "def\n"; }
    S(const S& other) { cout << "copy\n"; }
};

S f()
{
    S ret;
    return ret;
}

int main()
{
    S s = f();
    return 0;
}
```

??

Special Methods – RVO

```
struct S
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    S() { cout << "def\n"; }
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```
S f()
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    S ret;
    return ret;
}
```

```
int main()
{
    S s = f();
    return 0;
}
```

def g++ 4.4.3, no options

def
copy

def
copy
copy

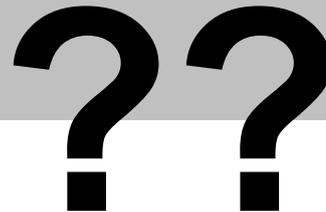


g++ -fno-elide-constructors

Special Methods – Copy Elision

```
struct S
{
    S(int i) { cout << i << "\n"; }
    S(const S&) { cout << "copy\n"; }
};

int main()
{
    S s1(1);
    S s2 = 2;
    return 0;
}
```



Special Methods – Copy Elision

```
struct S
{
    S(int i) { cout << i << "\n"; }
    S(const S&) { cout << "copy\n"; }
};
```

```
int main()
{
    S s1(1);
    S s2 = 2;
    return 0;
}
```

1

2



g++ 4.4.3, no options

1

2

copy



g++ -fno-elide-constructors

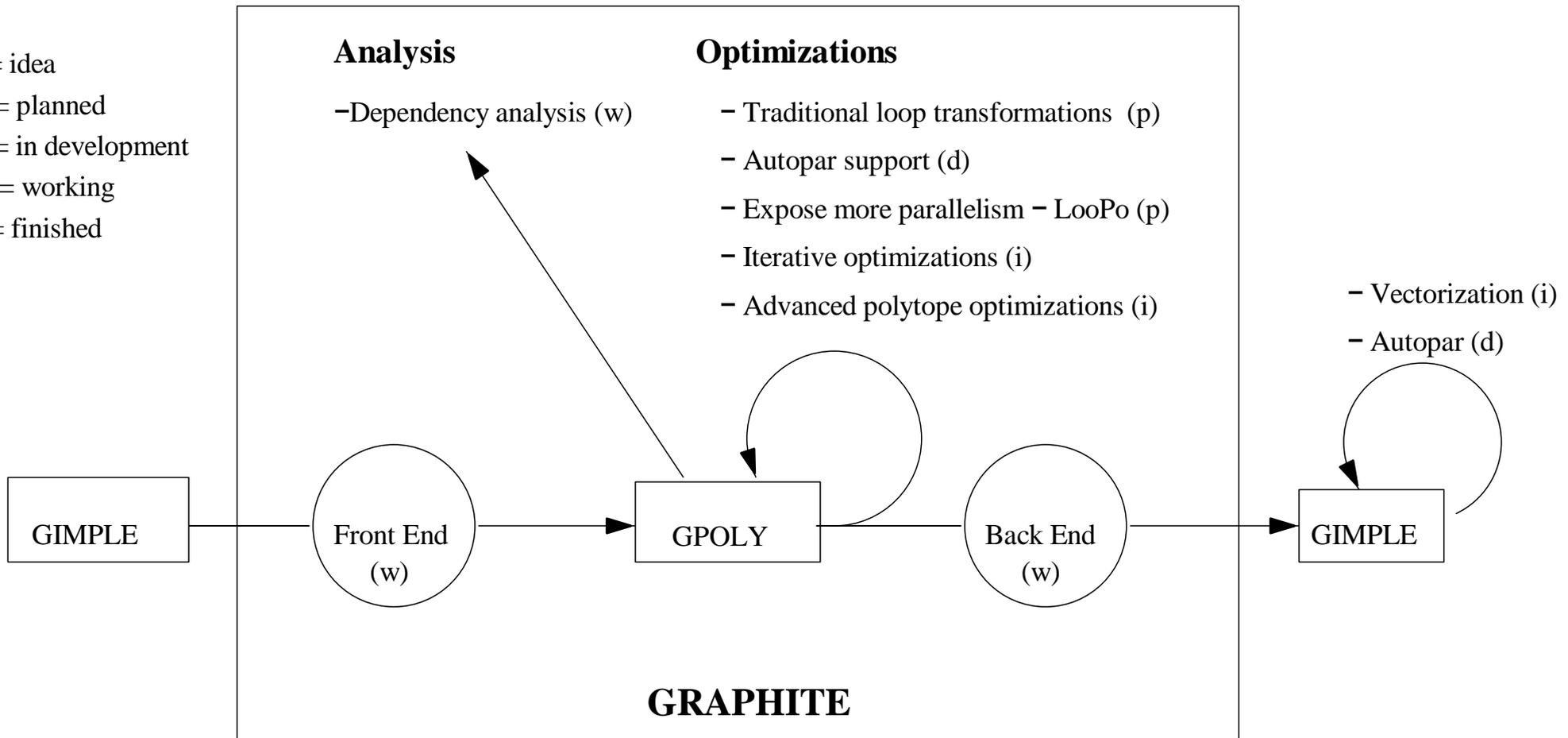
Optimizaciones Middle-end

- Transformaciones de loops
 - -funroll-loops
 - -fprefetch-loop-arrays
 - Graphite

Optimizaciones Middlend

Ej: Graphite

i = idea
p = planned
d = in development
w = working
f = finished



Graphite: Gimple Represented as Polyhedra

Graphite is a framework for high-level memory optimizations using the polyhedral model.

Graphite branch has been merged in August 2008 into trunk.

`-floop-parallelize-all`

Use the Graphite data dependence analysis to identify loops that can be parallelized. Parallelize all the loops that can be analyzed to not contain loop carried dependences without checking that it is profitable to parallelize the loops.

Optimizaciones Back-end

- Backend:
 - RTL tree + RTL language + RTL Engine + RTL Compiler
- [Middle-End] ----RTL tree→ [Back-end: Treewalker]
- Register Transfer Language:
 - RTL: Lisp-like
 - Puede embeber código C
 - Puede invocar funciones en C
 - Describe: conj. REGLAS
- Cada RTL Rule:
 - Tree Pattern
 - Condición (incluye código C)
 - Output assembly
 - Atributos (arch-specific), puntaje
- Score function

Back-end impl.

- Fuentes involucrados:
 - *arch.md*: Código RTL (MD: Machine Descrip.)
 - *arch.h*: Definiciones (macros) describen arch (qué reg. hay, para qué se usan, ...)
 - *arch.c*: Helper functions en C + *arch_rtx_costs()*

RTL

- RTL Model:
 - Máquina abstracta
 - Cantidad infinita de registros (pseudo-registers)
 - Realiza varias pasadas sobre sí mismo.
- La PRIMER OPTIMIZACION:
 - Al buildear GCC :-) → reglas para optimizar reglas

RTL

- Instruction scheduling pass
 - Tiene en cuenta el instruction decoder (ciclos), uOps
 - ***“Assembly/Compiler Coding Rule 25. (M impact, M generality)***
Avoid putting explicit references to ESP in a sequence of stack operations (POP, PUSH, CALL, RET).”
- Instruction selection pass
 - Tiene en cuenta subcores, pipeline
- Reloading Pass
 - *Asigna pseudo-registros en registros reales*
 - Cuando se queda sin registros reales: baja a memoria.
 - Una regla: cómo hacer reloading (insns, etc.)

- **Instruc**

- Tiene

- "Asse

- Avoid p
(POP, F

- **Instruc**

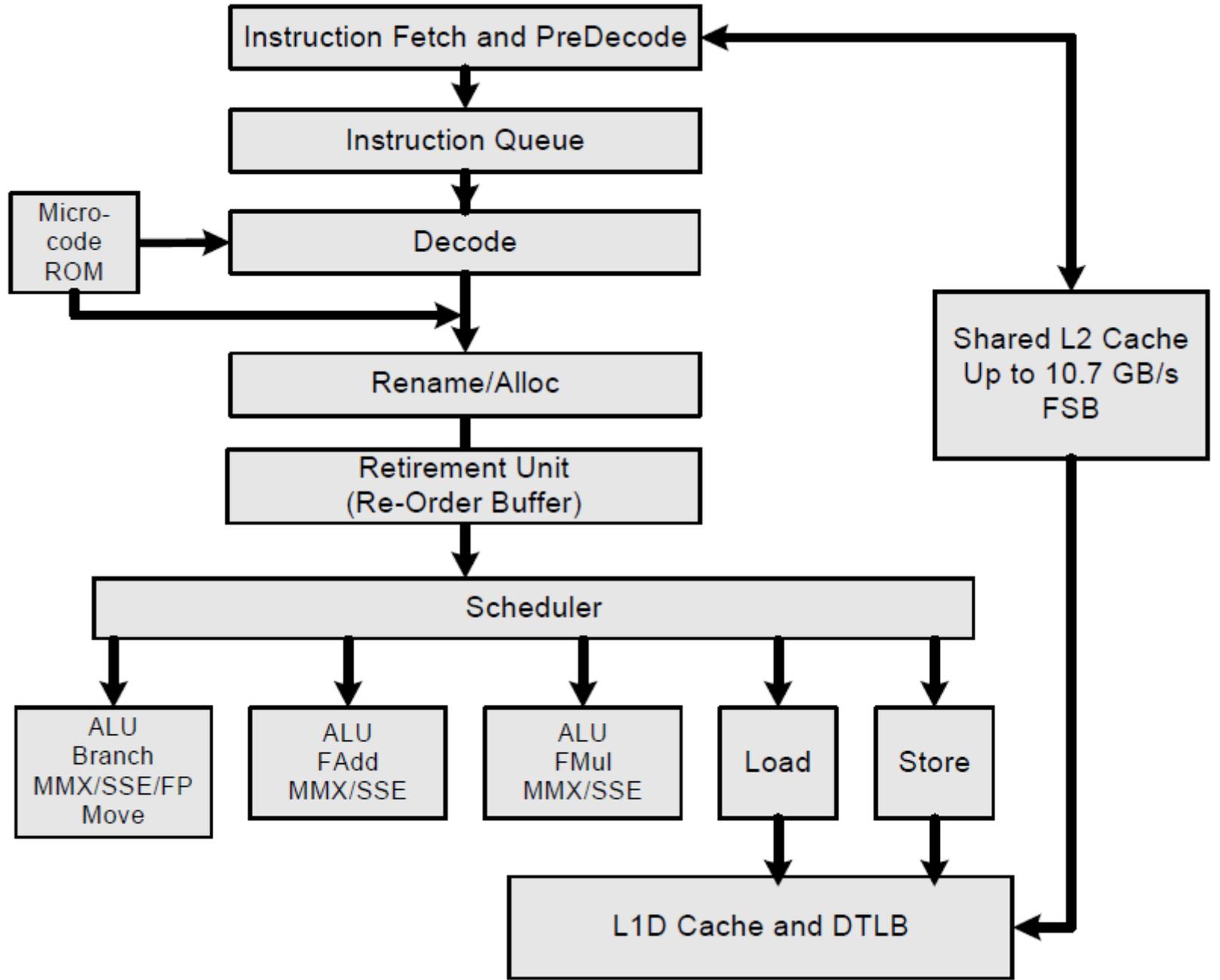
- Tiene

- **Reload**

- Asigr

- Cu

- Una regia. como hacer reloading (instns, etc.)



RTL

- Instruction scheduling pass
 - Tiene en cuenta el instruction decoder (ciclos), uOps
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 - Cuando se queda sin registros reales: baja a memoria.
 - Una regla: cómo hacer reloading (insns, etc.)

RTL

- Hojas RTL iniciales tienen **tipo específico** (cross-platform)
 - addi, addi3, subi, imul, ...
- Con cada pasada
 - Aplican reglas
 - Se van transformando/expandiendo/uniendo en instrucciones

RTL: Ejemplo real

- MADD: Multiply And Add
- En C:
 $a = a + b * c$
- Input RTL simplificado:
 - NODO_asignación { a, NODO_suma { a , NODO_multiplicacion { b, c } } }
- RTL rule:

```
(define_insn "*madds"  
  [(set (match_operand:ANYF 0 "register_operand" "=f")  
        (plus:ANYF (mult:ANYF (match_operand:ANYF 1 "register_operand" "f")  
                               (match_operand:ANYF 2 "register_operand" "f"))  
                  (match_operand:ANYF 3 "register_operand" "f")))]  
  "ISA_HAS_FP4 && TARGET_FUSED_MADD"  
  "madd.s %0,%3,%1,%2"  
  [(set_attr "type" "fmadd")  
   (set_attr "mode" "<UNITMODE>")])])
```

RTL: Ejemplo real

- MADD: Multiply And Add

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a = a + b*c

- Input RTL simplificado:

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                  (match_operand:ANYF 3 "register_operand" "f")))]
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  "madd.s %0,%3,%1,%2"
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```

Nombre
(* → debuggig)

RTL: Ejemplo real

- MADD: Multiply And Add

- En C:

a = a + b*c

- Input RTL simplificado:

- NODO_asignación { a, NODO_suma { a, NODO_multiplicacion { b, c } } }

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```
[(set (match_operand:ANYF 0 "register_operand" "=f")
      (plus:ANYF (mult:ANYF (match_operand:ANYF 1 "register_operand" "f")
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"ISA_HAS_FP4 && TARGET_FUSED_MADD"

"madd.s %0,%3,%1,%2"

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```

```
(set_attr "mode" "<UNITMODE>")])]
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- MADD: Multiply And Add

- En C:

```
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```

- Input RTL simplificado:

```
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```

- RTL rule:

```
(define_insn "*madds"
```

```
  [(set (match_operand:ANYF 0 "register_operand" "=f")
        (plus:ANYF (mult:ANYF (match_operand:ANYF 1 "register_operand" "f")
                               (match_operand:ANYF 2 "register_operand" "f"))
                  (match_operand:ANYF 3 "register_operand" "f")))]
```

```
"ISA_HAS_FP4 && TARGET_FUSED_MADD"
```

```
"madd.s %0,%3,%1,%2"
```

```
[(set_attr "type" "fmadd")
```

```
(set_attr "mode" "<UNITMODE>")]]
```

register_operand()
(mips.c)

RTL: Ejemplo real

- MADD: Multiply And Add
- En C:
 `a = a + b*c`
- Input RTL simplificado:
 - `NODO_asignación { a, NODO_suma { a , NODO_multiplicacion { b, c } } }`
- RTL rule:

```
(define_insn "*madds"  
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                               (match_operand:ANYF 2 "register_operand" "f"))  
                  (match_operand:ANYF 3 "register_operand" "f")))  
    "ISA_HAS_FP4 && TARGET_FUSED_MADD"  
    "madd.s %0,%3,%1,%2"  
  [(set_attr "type" "fmadd")  
   (set_attr "mode" "<UNITMODE>")])])
```

**#define ISA_HAS_FP4 ...
(mips.h)**

RTL: Ejemplo real

- MADD: Multiply And Add

- En C:

```
a = a + b*c
```

- Input RTL simplificado:

```
- NODO_asignación { a, NODO_suma { a , NODO_multiplicacion { b, c } } }
```

- RTL rule:

```
(define_insn "*madds"  
  [(set (match_operand:ANYF 0 "register_operand" "=f")  
        (plus:ANYF (mult:ANYF (match_operand:ANYF 1 "register_operand" "f")  
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  "madd.s %0,%3,%1,%2"  
  [(set_attr "type" "fmadd")  
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```

RTL: Ejemplo real

- MADD: Multiply And Add

- En C:

```
a = a + b*c
```

- Input RTL simplificado:

```
- NODO_asignación { a, NODO_suma { a , NODO_multiplicacion { b, c } } }
```

- RTL rule:

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  "madd.s %0,%3,%1,%2"  
  [(set_attr "type" "fmadd")  
   (set_attr "mode" "<UNITMODE>")])
```

**Scheduling rules,
cost**

RTL: Ejemplo real 2

- <http://gcc.gnu.org/ml/gcc-patches/2009-07/msg00009.html>

[PATCH] Workaround for Janus 2CC core errata

- *From:* Daniel Gutson <dgutson at codesourcery dot com>
- *To:* gcc-patches at gcc dot gnu dot org
- *Date:* Wed, 01 Jul 2009 03:26:14 -0300
- *Subject:* **[PATCH] Workaround for Janus 2CC core errata**

The attached patch implements a work-around for an errata in the Janus 2CC core.

I tested this by running the gcc test suite, including the two new test cases I added in this patch.

If accepted, please commit it for me since I don't have write access.

Thanks,
Daniel.

WPO/LTO vs Linker (gcc vs binutils)

- Contexto: ld (binutils) → lenguaje inespecífico
- OK para C, Fortran (70's)
 - Ej: init. variables globales: .bss
- NO OK para C++ (90's)
 - Ej: ctor/dtor objetos globales

collect2

- GCC's fake linker
 - Responsible for collecting calls for ctors & dtors in the .o files
 - Creates tables
 - Creates `__do_global_ctors_aux()` and the `__static_initialization_and_destruction_0()`
 - The latter: one per .o having globals
 - Has a call to `__cxa_atexit`
- Procedure:
 - Links the program once with the *real linker*
 - Collects data from binary (using nm)
 - Generates and compiles a temporary .c file
 - Re links everything with the new .o file with the *real linker*
- Verbose: add -debug to the linker (from gcc: -Wl,-debug)

Uso de collect2 para optimizar

- “Ve todo”
- Sacar ventaja si se usa *gold*: -fuse-linker-plugin
- IPA: Inter-Procedure analysis: callgraph
 - LGEN:
 - genera summary
 - graba summary
 - WPA
 - lee summary
 - “execute” (calcula opt.)
 - escribe opt. data
 - LTRANS:
 - lee opt. Data
 - Transforma (“transform”)

IPA: algunas optimizaciones

- Puede detectar funciones no usadas
- Permite cross-inlining en loops ajustados
- Propagation: puede saber valores de argumentos
- Puede saber valores de punteros
 - a objetos
 - a funciones

=> ahorra indirecciones.

Preguntas?

No, gracias.

Gracias!

Daniel.gutson@fudepan.org.ar