

Peak Performance for an Application in CUDA

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Outline

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Heat Conduction in a Plate

Also known as Laplace Differential Equation in two dimensions.

$$\nabla^2 T = \frac{\partial^2 T}{\partial x^2} + \frac{\partial^2 T}{\partial y^2} = 0$$

It can be solved iteratively using:

- Jacobi
- Gauss-Seidel
- Red-Black Gauss-Seidel

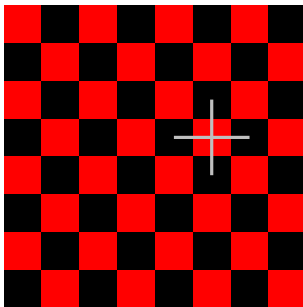
Red-Black Gauss-Seidel

A highly parallel, memory efficient method

Description

1st pass: update red cells from n to $n + 1$ using blacks in n .

2nd pass: update black cells from n to $n + 1$ using reds in $n + 1$.



Red-Black Gauss-Seidel Iteration

Red

For each **red** (i, j) , that is $(i + j) \bmod 2 = 0$ do:

$$M(i, j) \doteq \frac{M(i-1, j) + M(i, j-1) + M(i+1, j) + M(i, j+1)}{4}$$

Black

For each **black** (i, j) , that is $(i + j) \bmod 2 = 1$ do:

$$M(i, j) \doteq \frac{M(i-1, j) + M(i, j-1) + M(i+1, j) + M(i, j+1)}{4}$$

(stencil average)

Borders are **not** updated since they are our *contour conditions*.

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Serial_Checkboard

Remarks:

- Massive use of asserts
... they can be disabled with `-DNDEBUG` define.
- Border of fire of quadratic temperature.
Maximum in the center.
- Write simple PPMs image files as Pre/Post state.
- Measure **walltime** with microsecond precision.
- Print stats: **GFlops** and **GBps**.

Serial_Checkboard, results

```
gcc -std=gnu99 -O3 -Wall -DNDEBUG -DMAX_ITERATIONS=1000 -c -o  
gcc -std=gnu99 -O3 -Wall -DNDEBUG -DMAX_ITERATIONS=1000 -c -o  
gcc heat.o ../../common/common.o -o heat -std=gnu99 -O3 -Wall -D  
./heat
```

Secs: 12.777368

GBps: 1.641302

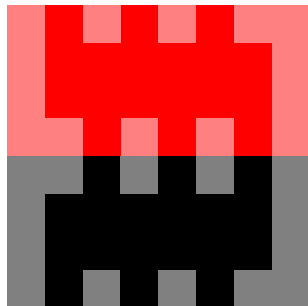
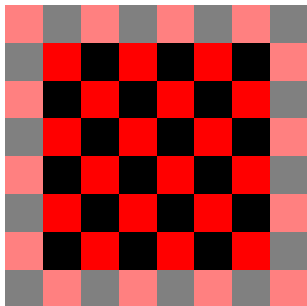
GFlops: 0.328260

Disclaimer

Our **GBps**, our **GFlops**, perhaps nothing to do with the real x86 architecture, but it serves for comparison.

Serial_Checkboard_Packed

$$(i, j) \rightarrow \left(\frac{(i+j) \bmod 2 \times N + i}{2}, j \right)$$



Serial_Checkboard_Packed, hoping less cache misses

With this transformation there are **three reads in a row**.

Thanks to cache *spatial locality* they implement **prefetching** bringing a whole **cache line**.

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Running it

```
./heat
```

```
Secs: 24.516243
```

```
GBps: 0.855413
```

```
GFlops: 0.171083
```

Not so useful.

Serial_Checkboard_Packed, less cache misses?

Using tools/perf/perf to **quantify** the executions.

```
Serial_Checkboard$ perf stat ./heat
Secs: 13.877461
GBps: 1.511193
GFlops: 0.302239
```

Performance counter stats for './heat':

```
29989928783 cycles
17154205308 instructions
133880188 cache-references
3208543 cache-misses
```

14.008524450 seconds time elapsed

```
Serial_Checkboard_Packed$ perf stat ./heat
Secs: 22.524494
GBps: 0.931054
GFlops: 0.186211
```

Performance counter stats for './heat':

```
48381072692 cycles
26476442265 instructions
138237703 cache-references
4295066 cache-misses
```

22.827983029 seconds time elapsed

It seems

- Double **i**nstruction **c**ount.
- No cache hits payoff.

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Peak performance for Red-Black Gauss-Seidel

- Memory-bandwidth limited problem.
- Measure memory bandwidth for the architecture.
- Use the knowledge of the underlying architecture to reach such value.

We are going to measure the BW using **SAXPY**.

$$\vec{y} = \alpha \vec{x} + \vec{y}$$

SAXPY/Serial

Architecture

Intel Core 2 Duo E8400, 2GB DDR3 1333MHz
(*January 2008*)

SAXPY measurment

For a vector of 2^{24} floats.

Secs: 0.030218

GBps: 6.662472

GFlops: 1.110412

Wikipedia informs **10.67 GBps**.

SAXPY/CUDA

Architecture

NVIDIA GTX 280 (GT200 arch), 1GB GDDR3, 30*8=240 cores
(*June 2008*)

SAXPY measurment

For a vector of 2^{24} floats, blocks of 512 threads.

Secs: 0.001747

GBps: 115.241323

GFlops: 19.206887

Wikipedia informs **141.70 GBps**.

Speedup

For a memory bandwidth intensive application:

17x

We aim to **peak**

in our app.

115GBps

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code inspection

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We emphasize

- All red(black) 512×1024 threads are divided in
 - A **block** of 16×16 threads.
 - A **grid** of 32×64 blocks.
- First map thread position to matrix element.
- Then compute the average.

We are (still) surprised with

- Extra fine granularity for concurrency.
- Pure C plus minor extensions.

CUDA_Checkboard, results

Running it

Secs: 0.403626

GBps: 51.957803

GFlops: 10.391561

Speedup

$$\frac{\textit{Serial_Checkboard}}{\textit{CUDA_Checkboard}} = \frac{12.777368}{0.403626} \geq 31x$$

Correctness

```
$ diff Serial_Checkboard/image_after.ppm CUDA_Checkboard/image_after.ppm  
Binary files Serial_Checkboard/... and CUDA_Checkboard/... differ
```

Not so...

CUDA_Checkboard_Packed

code inspection

CUDA_Checkboard_Packed

code inspection

Remarks

- Avoid branches using arithmetic.
- The only `if (...) { ... }` is for border detection.
- **Divergent branches** makes **warps** slow.

We hope for the best.

CUDA_Checkboard_Packed, results

Running it

Secs: 0.266050

GBps: 78.825484

GFlops: 15.765097

Speedup

$$\frac{\textit{Serial_Checkboard}}{\textit{CUDA_Checkboard}} = \frac{12.777368}{0.266050} \geq 48x$$

Do some **performance measurments** to find out **why**.

CUDA_Checkboard{ _Packed }, profiling

CUDA_Checkboard

```
# CUDA_PROFILE_LOG_VERSION 1.5
# CUDA_DEVICE 0 GeForce GTX 280
# TIMESTAMPFACTOR fe5c7d8438c7794
method,gputime,cputime,registerPerThread,occupancy,gld_32b,gld_64b,gld_128b,divergent_branch
method=[ memcpyHtoD ] gputime=[ 1535.264 ] cputime=[ 1865.000 ]
method=[ _Z8heat_redPf ] gputime=[ 200.032 ] cputime=[ 215.000 ] registerPerThread=[ 6 ]
    occupancy=[ 1.000 ] gld_32b=[ 3162 ] gld_64b=[ 22914 ] gld_128b=[ 3274 ] divergent_branch=[ 28 ]
method=[ _Z10heat_blackPf ] gputime=[ 203.104 ] cputime=[ 211.000 ] registerPerThread=[ 6 ]
    occupancy=[ 1.000 ] gld_32b=[ 3178 ] gld_64b=[ 22915 ] gld_128b=[ 3273 ] divergent_branch=[ 20 ]
```

CUDA_Checkboard_Packed

```
# CUDA_PROFILE_LOG_VERSION 1.5
# CUDA_DEVICE 0 GeForce GTX 280
# TIMESTAMPFACTOR fe5c7d92dd8961c
method,gputime,cputime,registerPerThread,occupancy,gld_32b,gld_64b,gld_128b,divergent_branch
method=[ memcpyHtoD ] gputime=[ 381.472 ] cputime=[ 1111.000 ]
method=[ memcpyHtoD ] gputime=[ 381.440 ] cputime=[ 1030.000 ]
method=[ _Z8heat_redPfS_ ] gputime=[ 134.752 ] cputime=[ 149.000 ] registerPerThread=[ 6 ]
    occupancy=[ 1.000 ] gld_32b=[ 3162 ] gld_64b=[ 13113 ] gld_128b=[ 3274 ] divergent_branch=[ 29 ]
method=[ _Z10heat_blackPfS_ ] gputime=[ 136.992 ] cputime=[ 146.000 ] registerPerThread=[ 6 ]
    occupancy=[ 1.000 ] gld_32b=[ 3178 ] gld_64b=[ 13114 ] gld_128b=[ 3273 ] divergent_branch=[ 20 ]
```

Thanks to **hardware performance counters**.

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Lessons learnt

- Avoid branches using arithmetic boolean-encoding tricks:
 $i+(2*(j\%2)-1)$
- Avoid (non) optimizations that hinder readability:
 $i!=(N-1)/2*(j\&1)$ identical to $i!=(N-1)/2*(j\%2)$
 $i>>=1$ identical to $i/=2$.
- `nvcc` is a **modern** compiler.
- Disconnect profiling: `export CUDA_PROFILE=0`.
- Compile to PTX and briefly inspect the code.
We discovered $0.25f*(...)$ is way different to $0.25*(...)$
- Play with block size and grid size in powers of two.
- Use all the (scarce) runtime error detection mechanisms:
`CUDA_SAFE_CALL(...)`, `CUT_CHECK_ERROR(...)`.
- Profile, measure, experiment, propose explaining hypothesis, etc.

Conclusions

We obtained **67%** of peak memory bandwidth, and a considerable speedup: **57x**.

The code is still **readable** for humans.

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Thanks!

Tomorrow we are going to deal with summarizations of the grid.
We need to know *when to stop*.

The Attic

Spare time?

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Spare time? Ok, let's show the little **weirdos** of the attic.

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What if we just program the **trivial** Jacobi iteration?

- Double matrix (swap old&new each turn).
- Fully parallel.

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- Double matrix (swap old&new each turn).
- Fully parallel.

Serial_Trivial

Secs: 25.292405

GBps: 0.829163

GFlops: 0.165833

Too bad for the serial case.

The Attic, the freak: `CUDA_Trivial`

code inspection

The Attic, the freak: `CUDA_Trivial`

code inspection

Running it

Secs: 0.218067

GBps: 96.170076

GFlops: 19.234015

Speedup

$$\frac{\textit{Serial_Checkboard}}{\textit{CUDA_Trivial}} = \frac{12.777368}{0.218076} \geq 58x$$

Amazing: zero effort, big gains, near peak bandwidth.

The Attic, the freak, the counters

CUDA_Checkboard_Packed

```
# CUDA_PROFILE_LOG_VERSION 1.5
# CUDA_DEVICE 0 GeForce GTX 280
# TIMESTAMPFACOR fe5c7d92dd8961c
method,gputime,cputime,registerPerThread,occupancy,gld_32b,gld_64b,gld_128b,divergent_branch
method=[ memcpyHtoD ] gputime=[ 381.472 ] cputime=[ 1111.000 ]
method=[ memcpyHtoD ] gputime=[ 381.440 ] cputime=[ 1030.000 ]
method=[ _Z8heat_redPfs_ ] gputime=[ 134.752 ] cputime=[ 149.000 ] registerPerThread=[ 6 ]
    occupancy=[ 1.000 ] gld_32b=[ 3162 ] gld_64b=[ 13113 ] gld_128b=[ 3274 ] divergent_branch=[ 29 ]
method=[ _Z10heat_blackPfs_ ] gputime=[ 136.992 ] cputime=[ 146.000 ] registerPerThread=[ 6 ]
    occupancy=[ 1.000 ] gld_32b=[ 3178 ] gld_64b=[ 13114 ] gld_128b=[ 3273 ] divergent_branch=[ 20 ]
```

CUDA_Trivial

```
# CUDA_PROFILE_LOG_VERSION 1.5
# CUDA_DEVICE 0 GeForce GTX 280
# TIMESTAMPFACOR fe5c7db01f0ad18
method,gputime,cputime,registerPerThread,occupancy,gld_32b,gld_64b,gld_128b,divergent_branch
method=[ memcpyHtoD ] gputime=[ 1535.424 ] cputime=[ 1872.000 ]
method=[ memcpyHtoD ] gputime=[ 1548.256 ] cputime=[ 1934.000 ]
method=[ _Z4heatPfs_ ] gputime=[ 220.224 ] cputime=[ 238.000 ] registerPerThread=[ 6 ]
    occupancy=[ 1.000 ] gld_32b=[ 6340 ] gld_64b=[ 19641 ] gld_128b=[ 6547 ] divergent_branch=[ 44 ]
```

There is no apparent reason, everything is worse.

The Attic, the freak, the reasons

Reading from global memory is a hairy business:

- Alignment
- How warps access the memory in each instruction

We hope a different transformation (parallelogram?) can render better results.

The Attic, the freak using shared: `CUDA_Trivial_Shared`

code inspection

The Attic, the freak using shared: CUDA_Trivial_Shared

code inspection

Running it

Secs: 0.242293

GBps: 86.554378

GFlops: 17.310876

Profiling it

```
method,gputime,cputime,registerPerThread,occupancy,gld_32b,gld_64b,gld_128b,divergent_branch
method=[ memcpyHtoD ] gputime=[ 1535.200 ] cputime=[ 1872.000 ]
method=[ memcpyHtoD ] gputime=[ 1544.320 ] cputime=[ 1932.000 ]
method=[ _Z4heatPFS_ ] gputime=[ 244.096 ] cputime=[ 262.000 ] registerPerThread=[ 7 ]
occupancy=[ 1.000 ] gld_32b=[ 4441 ] gld_64b=[ 4680 ] gld_128b=[ 3989 ] divergent_branch=[ 1530 ]
```

Each cell is read **four** times. Although in shared (ten to a hundred times faster), the more miss-aligned reads, the logic, and **divergent branches** make it slower.

Finally, two experiments

- Completely simetrical problem, swap i, j and see.
- *TILE* – 2 seems awkward to memory alignmen, what if just *TILE*?

See you tomorrow.

Questions?