Dynamic Epistemic Logics of Introspection

Raul Fervari¹ and Fernando R. Velázquez-Quesada²

¹FaMAF,UNC and CONICET, Argentina ²ILLC, UvA, The Netherlands

Workshop DaLí, Brasília, Brazil September 2017

Motivation

- Epistemic Logics deal with the knowledge of agents.
- Not only about propositional facts (ontic knowledge), but with the knowledge about her own and/or other agents' knowledge (high-order knowledge).
- Two important notions in high-order knowledge: positive and negative introspection.
- Positive introspection: "if the agent knows something, then she knows that she knows it".
- Negative introspection: "if the agent does not know something, then she knows that she does not know it".

Positive introspection: "if the agent knows something, then she knows that she knows it".

- Positive introspection: "if the agent knows something, then she knows that she knows it".
- In epistemic logic, it is characterized by

 $\Box \varphi \to \Box \Box \varphi$

- Positive introspection: "if the agent knows something, then she knows that she knows it".
- In epistemic logic, it is characterized by

$\Box \varphi \to \Box \Box \varphi$

• At semantic level, the accessibility relation is transitive.

- Positive introspection: "if the agent knows something, then she knows that she knows it".
- In epistemic logic, it is characterized by

$\Box \varphi \to \Box \Box \varphi$

- At semantic level, the accessibility relation is transitive.
- Negative introspection: "if the agent does not know something, then she knows that she does not know it".

- Positive introspection: "if the agent knows something, then she knows that she knows it".
- In epistemic logic, it is characterized by

$\Box \varphi \to \Box \Box \varphi$

- At semantic level, the accessibility relation is transitive.
- Negative introspection: "if the agent does not know something, then she knows that she does not know it".
- In epistemic logic, it is characterized by

$$\neg \Box \varphi \to \Box \neg \Box \varphi$$

- Positive introspection: "if the agent knows something, then she knows that she knows it".
- In epistemic logic, it is characterized by

$\Box \varphi \to \Box \Box \varphi$

- At semantic level, the accessibility relation is transitive.
- Negative introspection: "if the agent does not know something, then she knows that she does not know it".
- In epistemic logic, it is characterized by

$$\neg \Box \varphi \to \Box \neg \Box \varphi$$

• At the semantic level, the accessibility relation is Euclidean.

 We discussed introspection as properties in the logic and in the underlying structures (relational models).

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?
- Properties as the eventual result of an action:

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?
- Properties as the eventual result of an action:
 - "Syntatic" inference steps based on sets of formulas

[Duc17,Ågo&Alech07,Jago09].

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?
- Properties as the eventual result of an action:
 - "Syntatic" inference steps based on sets of formulas

[Duc17,Ågo&Alech07,Jago09]. [Gross&VQ15].

Inference steps in awareness models

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?
- Properties as the eventual result of an action:
 - "Syntatic" inference steps based on sets of formulas

[Duc17, Ågo&Alech07, Jago09].

- Inference steps in awareness models
- Dynamics of evidence, deductive inference [VBen&Pacuit11,VQ13].

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?
- Properties as the eventual result of an action:
 - "Syntatic" inference steps based on sets of formulas

```
[Duc17,Ågo&Alech07,Jago09].
```

- Inference steps in awareness models
- Dynamics of evidence, deductive inference [VBen&Pacuit11,VQ13].
- Dynamic Epistemic Logic style.

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?
- Properties as the eventual result of an action:
 - "Syntatic" inference steps based on sets of formulas

[Duc17,Ågo&Alech07,Jago09].

- Inference steps in awareness models
- Dynamics of evidence, deductive inference [VBen&Pacuit11,VQ13].
- Dynamic Epistemic Logic style.
- Operations that change the model:
 - Actions for BR, preferences [vBen07,vBen&Liu07,Ghosh&VQ15].

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?
- Properties as the eventual result of an action:
 - "Syntatic" inference steps based on sets of formulas

[Duc17,Ågo&Alech07,Jago09].

- Inference steps in awareness models
- Dynamics of evidence, deductive inference [VBen&Pacuit11,VQ13].
- Dynamic Epistemic Logic style.
- Operations that change the model:
 - Actions for BR, preferences [vBen07,vBen&Liu07,Ghosh&VQ15].
 - Sabotage logic

[vBen05].

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?
- Properties as the eventual result of an action:
 - "Syntatic" inference steps based on sets of formulas

[Duc17,Ågo&Alech07,Jago09].

- Inference steps in awareness models
- Dynamics of evidence, deductive inference [VBen&Pacuit11,VQ13].
- Dynamic Epistemic Logic style.

Sabotage logic

- Operations that change the model:
 - Actions for BR, preferences [vBen07,vBen&Liu07,Ghosh&VQ15].
 - [vBen05].

[Gross&VQ15].

[AFH12,14,15,Ferv14].

Relation-changing modal logics

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?
- Properties as the eventual result of an action:
 - "Syntatic" inference steps based on sets of formulas

```
[Duc17,Ågo&Alech07,Jago09].
```

- Inference steps in awareness models
- Dynamics of evidence, deductive inference [VBen&Pacuit11,VQ13].
- Dynamic Epistemic Logic style.
- Operations that change the model:

Relation-changing modal logics

- Actions for BR, preferences [vBen07,vBen&Liu07,Ghosh&VQ15].
 - [vBen05].

[Gross&VQ15].

- [AFH12,14,15,Ferv14].
 - [Kooi&Renne11].

Arrow updates

Sabotage logic

- We discussed introspection as properties in the logic and in the underlying structures (relational models).
- What about considering it as properties to be achieved?
- Properties as the eventual result of an action:
 - "Syntatic" inference steps based on sets of formulas

```
[Duc17,Ågo&Alech07,Jago09].
```

- Inference steps in awareness models
- Dynamics of evidence, deductive inference [VBen&Pacuit11,VQ13].
- Dynamic Epistemic Logic style.
- Operations that change the model:
 - Actions for BR, preferences [vBen07,vBe
 - Sabotage logic
 - Relation-changing modal logics
 - Arrow updates
 - More....

[vBen07,vBen&Liu07,Ghosh&VQ15].

- [vBen05].
- [AFH12,14,15,Ferv14].
 - [Kooi&Renne11].

Our work

- We work in a single-agent framework.
- Our approach is provide operations that after being executed, +/introspection is achieved.
- Two kind of operations (both in a global sense):
 - General introspection.
 - Introspection with respect to a formula.
- Two possible approaches:
 - Pessimistic: the agent loses some knowledge.
 - Optimistic: the agent gains some knowledge.

▶ First idea: to make $\Box \varphi \rightarrow \Box \Box \varphi$ true, make the accessibility relation transitive.

- ▶ First idea: to make $\Box \varphi \rightarrow \Box \Box \varphi$ true, make the accessibility relation transitive.
- Let $M = \langle W, R, V \rangle$ be a relational model, $w \in W$:

 $M, w \Vdash \langle + \rangle \varphi$ iff $M^+, w \Vdash \varphi$

with $M^+ = \langle W, R^+, V \rangle$, where R^+ is the transitive closure of R.

- ▶ First idea: to make $\Box \varphi \rightarrow \Box \Box \varphi$ true, make the accessibility relation transitive.
- Let $M = \langle W, R, V \rangle$ be a relational model, $w \in W$:

 $M, w \Vdash \langle + \rangle \varphi$ iff $M^+, w \Vdash \varphi$

with $M^+ = \langle W, R^+, V \rangle$, where R^+ is the transitive closure of R.

• Self-duality: $[+] \varphi := \neg \langle + \rangle \neg \varphi$, is equivalent to $\langle + \rangle$.

- ▶ First idea: to make $\Box \varphi \rightarrow \Box \Box \varphi$ true, make the accessibility relation transitive.
- Let $M = \langle W, R, V \rangle$ be a relational model, $w \in W$:

 $M, w \Vdash \langle + \rangle \varphi$ iff $M^+, w \Vdash \varphi$

with $M^+ = \langle W, R^+, V \rangle$, where R^+ is the transitive closure of R.

• Self-duality: $[+] \varphi := \neg \langle + \rangle \neg \varphi$, is equivalent to $\langle + \rangle$.

Proposition

$$\Vdash [+] (\Box \varphi \to \Box \Box \varphi).$$

General Positive Instrospection - some properties

▶ Complete axiomatization (reduction axioms to K⁺):

 $\vdash {\bf (+)} \diamondsuit \varphi \leftrightarrow \oplus {\bf (+)} \varphi$

General Positive Instrospection - some properties

• Complete axiomatization (reduction axioms to K^+):

$$\vdash \langle + \rangle \diamondsuit \varphi \leftrightarrow \oplus \langle + \rangle \varphi$$
$$\vdash \langle + \rangle \oplus \varphi \leftrightarrow \oplus \langle + \rangle \varphi$$

General Positive Instrospection - some properties

• Complete axiomatization (reduction axioms to K^+):

$$\vdash \langle + \rangle \diamondsuit \varphi \leftrightarrow \oplus \langle + \rangle \varphi$$
$$\vdash \langle + \rangle \oplus \varphi \leftrightarrow \oplus \langle + \rangle \varphi$$

However, is not the expected behaviour:

Fact

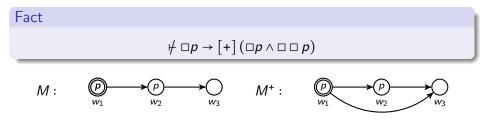
$$\not\vdash \Box p \rightarrow [+] (\Box p \land \Box \Box p)$$

Losing knowledge

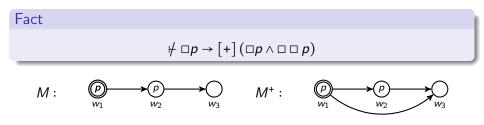
Fact

$$\not\vdash \Box p \rightarrow [+] (\Box p \land \Box \Box p)$$

Losing knowledge



Losing knowledge



Observation

We are getting introspection at cost of losing knowledge!

• A more intuitive approach: take the agent from knowing χ without knowing she knows it, to knowing χ and knowing she knows it.

- A more intuitive approach: take the agent from knowing χ without knowing she knows it, to knowing χ and knowing she knows it.
- If at (M, w) the agent knows a given χ without having full positive introspection about it, then:
 - 1. every world *R*-reachable from *w* in one step satisfies χ ,

- A more intuitive approach: take the agent from knowing χ without knowing she knows it, to knowing χ and knowing she knows it.
- If at (M, w) the agent knows a given χ without having full positive introspection about it, then:
 - 1. every world *R*-reachable from *w* in one step satisfies χ ,
 - 2. there is at least one world *R*-reachable from *w* (in two or more steps) where χ fails.

- A more intuitive approach: take the agent from knowing χ without knowing she knows it, to knowing χ and knowing she knows it.
- If at (M, w) the agent knows a given χ without having full positive introspection about it, then:
 - 1. every world *R*-reachable from *w* in one step satisfies χ ,
 - 2. there is at least one world *R*-reachable from *w* (in two or more steps) where χ fails.
- The operation should make the $\neg \chi$ -worlds inaccessible.

- A more intuitive approach: take the agent from knowing χ without knowing she knows it, to knowing χ and knowing she knows it.
- If at (M, w) the agent knows a given χ without having full positive introspection about it, then:
 - 1. every world *R*-reachable from *w* in one step satisfies χ ,
 - 2. there is at least one world *R*-reachable from *w* (in two or more steps) where χ fails.
- The operation should make the $\neg \chi$ -worlds inaccessible.
- We need to remove edges.

• Remove edges pointing to $\neg \chi$ -worlds.

- Remove edges pointing to $\neg \chi$ -worlds.
- Let $M = \langle W, R, V \rangle$ be a relational model; take $U \subseteq W$.

- Remove edges pointing to $\neg \chi$ -worlds.
- Let $M = \langle W, R, V \rangle$ be a relational model; take $U \subseteq W$.
- The *U*-disconnecting operation yields the model $M_{+U} = \langle W, R', V \rangle$, with $R' := R \setminus (U \times \overline{U})$ (for $\overline{U} := W \setminus U$).

- Remove edges pointing to $\neg \chi$ -worlds.
- Let $M = \langle W, R, V \rangle$ be a relational model; take $U \subseteq W$.
- The *U*-disconnecting operation yields the model $M_{+U} = \langle W, R', V \rangle$, with $R' := R \setminus (U \times \overline{U})$ (for $\overline{U} := W \setminus U$).
- Thus, this operation removes edges from worlds on U to worlds not in U.

- Remove edges pointing to $\neg \chi$ -worlds.
- Let $M = \langle W, R, V \rangle$ be a relational model; take $U \subseteq W$.
- The *U*-disconnecting operation yields the model $M_{+U} = \langle W, R', V \rangle$, with $R' := R \setminus (U \times \overline{U})$ (for $\overline{U} := W \setminus U$).
- Thus, this operation removes edges from worlds on U to worlds not in U.

$$(M,w) \Vdash \langle +'\chi \rangle \varphi \quad \text{iff} \quad (M_{+\llbracket \chi \rrbracket^M},w) \Vdash \varphi$$

- Remove edges pointing to $\neg \chi$ -worlds.
- Let $M = \langle W, R, V \rangle$ be a relational model; take $U \subseteq W$.
- The *U*-disconnecting operation yields the model $M_{+U} = \langle W, R', V \rangle$, with $R' := R \setminus (U \times \overline{U})$ (for $\overline{U} := W \setminus U$).
- ▶ Thus, this operation removes edges from worlds on U to worlds not in U.

$$(M,w) \Vdash \langle +'\chi \rangle \varphi \quad \text{iff} \quad (M_{+[\![\chi]\!]^M},w) \Vdash \varphi$$

• But... this operation can take place in any situation! (not only when the agent knows χ)

As we pointed out, we get the expected behaviour by including a pre-condition:

$$\langle +\chi \rangle \varphi \coloneqq \Box \chi \land \langle +'\chi \rangle \varphi.$$

As we pointed out, we get the expected behaviour by including a pre-condition:

$$\langle +\chi \rangle \varphi \coloneqq \Box \chi \land \langle +'\chi \rangle \varphi.$$

• The relation of the resulting model can be equivalently defined, using PDL notation, as $R := (?\neg\chi; R) \cup (R; ?\chi)$ (easy to axiomatize with reduction axioms).

As we pointed out, we get the expected behaviour by including a pre-condition:

$$\langle +\chi \rangle \varphi \coloneqq \Box \chi \land \langle +'\chi \rangle \varphi.$$

• The relation of the resulting model can be equivalently defined, using PDL notation, as $R \coloneqq (?\neg\chi; R) \cup (R; ?\chi)$ (easy to axiomatize with reduction axioms).

Proposition

$$\Vdash \langle +\chi \rangle \Box \varphi \iff \Box (\chi \land [+'\chi] \varphi).$$

As we pointed out, we get the expected behaviour by including a pre-condition:

$$\langle +\chi \rangle \varphi \coloneqq \Box \chi \land \langle +'\chi \rangle \varphi.$$

• The relation of the resulting model can be equivalently defined, using PDL notation, as $R \coloneqq (?\neg\chi; R) \cup (R; ?\chi)$ (easy to axiomatize with reduction axioms).

Proposition

$$\Vdash \langle +\chi \rangle \Box \varphi \iff \Box (\chi \land [+'\chi] \varphi).$$

"The agent can perform a particular positive introspection step for χ after which she will know φ iff she knows both χ and that, after the 'preconditionless' operation, φ will be the case."

As we pointed out, we get the expected behaviour by including a pre-condition:

$$\langle +\chi \rangle \varphi \coloneqq \Box \chi \land \langle +'\chi \rangle \varphi.$$

• The relation of the resulting model can be equivalently defined, using PDL notation, as $R := (?\neg\chi; R) \cup (R; ?\chi)$ (easy to axiomatize with reduction axioms).

Proposition

$$\Vdash \langle +\chi \rangle \Box \varphi \iff \Box (\chi \land [+'\chi] \varphi).$$

Replacing φ by $\Box \chi$ we get:

$$\vdash \langle +\chi \rangle \Box \Box \chi \leftrightarrow \Box (\chi \land [+'\chi] \Box \chi).$$

 PAL requires pre-condition χ to be true, but the introspection operation requires for χ to be known.

- PAL requires pre-condition \(\chi \) to be true, but the introspection operation requires for \(\chi \) to be known.
- Moorean phenomena: in PAL, after being truthfully announced, become false (e.g. p ∧ □¬p).

- PAL requires pre-condition \(\chi \) to be true, but the introspection operation requires for \(\chi \) to be known.
- Moorean phenomena: in PAL, after being truthfully announced, become false (e.g. p ∧ □¬p).
- Here we have:

Fact

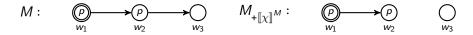
 $\not\models \Box \chi \to [+'\chi] \Box \chi, \text{ as a consequence } \not\models [+'\chi] \Box \chi.$

- PAL requires pre-condition \(\chi \) to be true, but the introspection operation requires for \(\chi \) to be known.
- Moorean phenomena: in PAL, after being truthfully announced, become false (e.g. p ∧ □¬p).
- Here we have:

Fact

$$\not\models \Box \chi \to [+'\chi] \Box \chi, \text{ as a consequence } \not\models [+'\chi] \Box \chi.$$

Take $\chi \coloneqq p \land \diamondsuit \diamondsuit \neg p$:

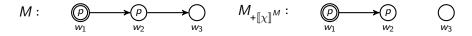


- PAL requires pre-condition \(\chi \) to be true, but the introspection operation requires for \(\chi \) to be known.
- Moorean phenomena: in PAL, after being truthfully announced, become false (e.g. p ∧ □¬p).
- Here we have:

Fact

$$\not\models \Box \chi \to [+'\chi] \Box \chi, \text{ as a consequence } \not\models [+'\chi] \Box \chi.$$

Take $\chi := p \land \diamondsuit \diamondsuit \neg p$:



Proposition

If $\Vdash \chi \rightarrow [+'\chi] \chi$, then after the operation the agent will have positive introspection about χ , i.e., $\Vdash \langle +\chi \rangle \Box \Box \chi \leftrightarrow \Box \chi$.

▶ First idea: to make $\neg \Box \varphi \rightarrow \Box \neg \Box \varphi$ true, make the accessibility relation Euclidean.

- First idea: to make $\neg \Box \varphi \rightarrow \Box \neg \Box \varphi$ true, make the accessibility relation Euclidean.
- Let M = ⟨W, R, V⟩ be a relational model, w ∈ W; M⁻ = ⟨W, R^E, V⟩ in which R^E is the Euclidean closure of R, that is,

$$R^E \coloneqq R \cup (\mathfrak{R} \circ (R \cup \mathfrak{R})^* \circ R).$$

- First idea: to make $\neg \Box \varphi \rightarrow \Box \neg \Box \varphi$ true, make the accessibility relation Euclidean.
- Let M = ⟨W, R, V⟩ be a relational model, w ∈ W; M⁻ = ⟨W, R^E, V⟩ in which R^E is the Euclidean closure of R, that is,

$$R^E \coloneqq R \cup (\mathfrak{R} \circ (R \cup \mathfrak{R})^* \circ R).$$

The associated modality:

 $(M, w) \Vdash \langle - \rangle \varphi$ iff $(M^-, w) \Vdash \varphi$.

- First idea: to make $\neg \Box \varphi \rightarrow \Box \neg \Box \varphi$ true, make the accessibility relation Euclidean.
- Let M = ⟨W, R, V⟩ be a relational model, w ∈ W; M⁻ = ⟨W, R^E, V⟩ in which R^E is the Euclidean closure of R, that is,

$$R^E \coloneqq R \cup (\mathfrak{R} \circ (R \cup \mathfrak{R})^* \circ R).$$

The associated modality:

 $(M, w) \Vdash \langle - \rangle \varphi$ iff $(M^-, w) \Vdash \varphi$.

Proposition

$$\Vdash [-] (\neg \Box \varphi \rightarrow \Box \neg \Box \varphi).$$

- First idea: to make $\neg \Box \varphi \rightarrow \Box \neg \Box \varphi$ true, make the accessibility relation Euclidean.
- Let M = ⟨W, R, V⟩ be a relational model, w ∈ W; M⁻ = ⟨W, R^E, V⟩ in which R^E is the Euclidean closure of R, that is,

$$R^E \coloneqq R \cup (\mathfrak{R} \circ (R \cup \mathfrak{R})^* \circ R).$$

The associated modality:

 $(M, w) \Vdash \langle - \rangle \varphi$ iff $(M^-, w) \Vdash \varphi$.

Proposition

$$\Vdash [-] (\neg \Box \varphi \to \Box \neg \Box \varphi).$$

Even more:

Proposition

If φ is propositional, then $\Vdash \neg \Box \varphi \rightarrow [-] (\neg \Box \varphi \land \Box \neg \Box \varphi)$.

As expected, it also has Moorean behaviour for arbitrary formulas:

- As expected, it also has Moorean behaviour for arbitrary formulas:

Proposition $|\not - \neg \Box \varphi \rightarrow [-] \Box \neg \Box \varphi.$

As expected, it also has Moorean behaviour for arbitrary formulas:



As expected, it also has Moorean behaviour for arbitrary formulas:

Axiomatizable in PDL (with test and converse).

 For the propositional case, the general operation already behaves as expected.

- For the propositional case, the general operation already behaves as expected.
- However, for uniformity we also define the particular version.

- For the propositional case, the general operation already behaves as expected.
- However, for uniformity we also define the particular version.
- Let M = ⟨W, R, V⟩ and U ⊆ W. The U-connecting operation gives the model M_{-U} = ⟨W, R', V⟩ with

$$R' \coloneqq R \cup (\mathfrak{R} \circ (R \cup \mathfrak{R})^* \circ R \circ \mathsf{Id}_U^M),$$

where $\operatorname{Id}_U^M \coloneqq \{(u, u) \mid u \in U\}.$

- For the propositional case, the general operation already behaves as expected.
- However, for uniformity we also define the particular version.
- Let M = ⟨W, R, V⟩ and U ⊆ W. The U-connecting operation gives the model M_{-U} = ⟨W, R', V⟩ with

$$R' \coloneqq R \cup (\mathfrak{R} \circ (R \cup \mathfrak{R})^* \circ R \circ \mathsf{Id}_U^M),$$

where
$$\operatorname{Id}_U^M \coloneqq \{(u, u) \mid u \in U\}.$$

We define a pre-conditionless operation:

$$(M, w) \Vdash \langle -'\chi \rangle \varphi \quad \text{iff} \quad (M_{-[\neg \chi]^M}, w) \Vdash \varphi.$$

- For the propositional case, the general operation already behaves as expected.
- However, for uniformity we also define the particular version.
- Let M = ⟨W, R, V⟩ and U ⊆ W. The U-connecting operation gives the model M_{-U} = ⟨W, R', V⟩ with

$$R' \coloneqq R \cup (\mathfrak{R} \circ (R \cup \mathfrak{R})^* \circ R \circ \mathsf{Id}_U^M),$$

where
$$\operatorname{Id}_U^M \coloneqq \{(u, u) \mid u \in U\}.$$

• We define a pre-conditionless operation:

$$(M, w) \Vdash \langle -'\chi \rangle \varphi \quad \text{iff} \quad (M_{-[\neg \chi]^M}, w) \Vdash \varphi.$$

And the corresponding modality:

$$\langle -\chi \rangle \varphi \coloneqq \neg \Box \chi \land \langle -'\chi \rangle \varphi.$$

 We define operations for achieving positive and negative introspection, both in general and also for a particular formula.

- We define operations for achieving positive and negative introspection, both in general and also for a particular formula.
- We studied their properties, and obtained complete axiomatizations (via reductions axioms into PDL).

- We define operations for achieving positive and negative introspection, both in general and also for a particular formula.
- We studied their properties, and obtained complete axiomatizations (via reductions axioms into PDL).
- For positive introspection: one operation adds edges, the other removes edges.

- We define operations for achieving positive and negative introspection, both in general and also for a particular formula.
- We studied their properties, and obtained complete axiomatizations (via reductions axioms into PDL).
- For positive introspection: one operation adds edges, the other removes edges.
- For negative introspection both operations adds edges.

Future Work

- Getting introspection in one level (more local): from $\Box p \land \neg \Box \Box p$ to $\Box p \land \Box \Box p \land \neg \Box \Box D p$
- We would like to explore this kind of operations in a multi-agent setting.
 - Public, private versions of these operations.
 - Reaching common knowledge.
- Connections with cognitive operations.